

### SOLUTION SET 3

**Problem 1** Let  $G$  be a graph with  $n$  vertices and  $e$  edges, so that every edge lies in  $m$  spanning trees. Consider the graph as made of identical resistors with resistance  $R$ . What is the effective resistance between two adjacent vertices?

Let  $u$  and  $v$  be two adjacent vertices. From class, the effective resistance between  $u$  and  $v$  is

$$R \frac{\#(\text{two tree forests separating } u \text{ and } v)}{\#(\text{spanning forests})}.$$

The numerator is the same as the number of spanning trees of  $G$  which contain the edge from  $u$  to  $v$ . So the effective resistance is the probability that a random tree will contain edge  $e$ .

Let  $t$  be the total number of spanning trees. Since each tree has  $n - 1$  edges, and each edge is in  $m$  trees, we have  $t(n - 1) = em$ . We want to compute  $m/t = (n - 1)/e$ . So the final answer is  $(n - 1)/e \times R$ .

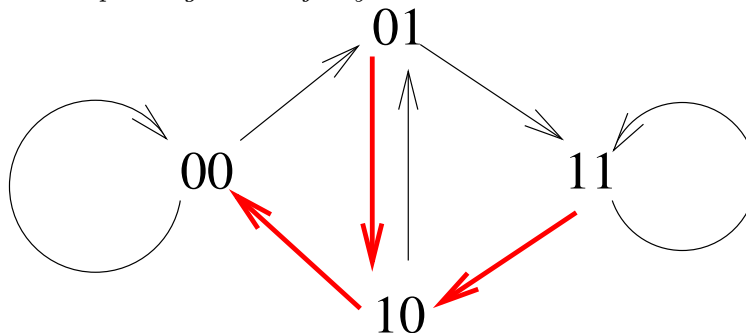
**Problem 2** Let  $n$  be a positive integer. Construct a binary string in a greedy manner as follows: Start with  $n - 1$  copies of 0. Add bits to the end as follows: Add 1 unless this would cause the same  $n$ -bit substring to occur twice. If we are forbidden to add 1, then add 0 if this will not cause an  $n$ -bit substring to reoccur. If both 1 and 0 would cause an  $n$ -bit substring to reoccur, then stop. For  $n = 3$ , this produces

0011101000.

In this problem, we will show that this procedure ends with  $n - 1$  zeroes and that gluing those zeroes to the initial zeroes produces a de Bruijn cycle.

Let  $D_n$  be the directed graph whose vertices are binary strings of length  $n - 1$  and where there is an edge from  $b_1b_2 \cdots b_{n-1}$  to  $b_2b_3 \cdots b_{n-1}b_n$ . (This is the graph we used in class to prove de Bruijn sequences exist.) Let  $T$  be the subgraph consisting of the edges  $b_1b_2 \cdots b_{n-1} \rightarrow b_2 \cdots b_{n-1}0$  for every binary string  $b_1b_2 \cdots b_{n-1}$  other than  $00 \cdots 0$ .

(a) For  $n = 3$ , draw the graph  $D_3$  and the subgraph  $T$ . You should see that  $T$  is a rooted subtree of  $D_3$ ; construct the corresponding de Bruijn cycle.



(b) Prove in general that  $T$  is a rooted subtree of  $D_n$ .

We will show that, for every vertex  $v$  of  $T$ , there is a directed path from  $v$  to  $00 \cdots 0$ . This shows that there cannot be any directed cycles, so  $T$  is rooted at  $00 \cdots 0$ . This is simple: Let  $v = v_1v_2 \cdots v_{n-1}$ . Then

$$(v_1v_2 \cdots v_{n-1}) \rightarrow (v_2 \cdots v_{n-1}0) \rightarrow (v_3 \cdots v_{n-1}00) \rightarrow \cdots \rightarrow (v_{n-1}0 \cdots 0) \rightarrow (00 \cdots 0)$$

is a directed path in  $T$ .

(c) Prove that the Eulerian walk of  $D_n$  corresponding to  $T$  (by the BEST algorithm) gives the de Bruijn sequence constructed by the greedy procedure above.

In the BEST algorithm, we first leave a given vertex using the edge not in  $T$  and then, the next time, we take the edge from  $T$ . (If the out-degree were greater than 2, we would also need a way to choose between the various out-going edges not in  $T$ , but the out-degree is 2.) The edge of  $T$

pointing out of  $v_1v_2 \cdots v_{n-1}$  is to  $v_1 \cdots v_{n-1}0$ , so this algorithm says to choose 1 before choosing 0, precisely as the greedy algorithm specifies.

**Problem 3** Let  $G$  be a directed graph with  $n$  vertices  $v_1, v_2, \dots, v_n$  arranged around a circle.  $G$  has  $4n$  edges: For each vertex, there are two edges to its clockwise neighbor and two to its counter-clockwise neighbor.

(a) Find the number of rooted spanning trees of  $G$ . (Just thinking about the problem is probably easier than using the matrix-tree theorem.)

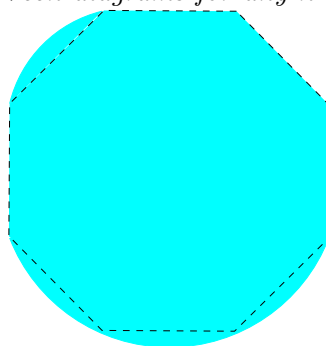
We must choose precisely one of the pairs  $(v_i, v_{i+1})$  not to be joined by an edge. We must choose exactly one of the vertices to be the root. Finally, for each of the  $n - 1$  pairs  $(v_j, v_{j+1})$  which are joined by an edge, we have 2 choices of which edge to use. This gives  $n^2 2^{n-1}$  trees total. If we specify a particular root, there are  $n 2^{n-1}$ .

(b) Find the number of Eulerian walks in  $G$ .

The number of Eulerian walks starting (and hence ending) at  $v_i$  will be  $(n 2^{n-1}) \times 4! \times (3!)^{n-1}$ . So the total number is  $n^2 \times (12)^{n-1} \times 24$ .

**Problem 4** Let  $n$  be a positive integer. Let  $X_1, X_2, \dots, X_n$  be subsets of the plane  $\mathbb{R}^2$ . For any subset  $I$  of  $\{1, 2, \dots, n\}$ , let  $Y_I = (\bigcap_{i \in I} X_i) \cap (\bigcap_{j \notin I} X_j^C)$ , where  $X_j^C$  is  $\mathbb{R}^2 \setminus X_j$ . We will define  $X_1, \dots, X_n$  to be a generalized Venn diagram if all the  $X_i$ 's are convex and all the  $Y_I$ 's are nonempty.

Show that there exist generalized Venn diagrams for any  $n$ .



Let  $P$  be a regular  $2^n$  inscribed in a circle. Label the  $2^n$  wedges between  $P$  and the circle with a binary de Bruijn cycle. Let  $X_1$  be the union of  $P$  and the wedges which are labeled 1. Let  $X_k$ , for  $2 \leq k \leq n$ , be the rotation of  $X_1$  by  $2\pi k/2^n$ . As can be seen from the figure,  $X_k$  is convex. For  $I$  other than  $\emptyset$  or  $\{1, 2, \dots, n\}$ , the intersection  $Y_I$  is a single wedge.  $Y_\emptyset$  is the exterior of the circle, plus one more wedge. Finally,  $Y_{\{1, 2, \dots, n\}}$  is the union of  $P$  and one wedge.