# NOTES FOR NOVEMBER 19, 2012: INTERACTION BETWEEN $\mathfrak{gl}_2$ STRINGS

### AARON PRIBADI

### 1. From last class ( $\mathfrak{gl}_2$ representations)

The Lie algebra  $\mathfrak{gl}_n$  consists of the  $n \times n$  matrices with bracket  $[,]: \mathfrak{gl}_n \times \mathfrak{gl}_n \to \mathfrak{gl}_n$  given by the matrix commutator [A, B] = AB - BA. It has a basis of elementary matrices  $E_{ij}$  with a single 1 in *i*th row and *j*th column, and 0's everywhere else.

Define

$$h_k = E_{k\,k}$$
  $e_k = E_{k\,(k+1)}$   $f_k = E_{(k+1)\,k}$ 

In particular,

$$h_1 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$
  $h_2 = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$   $e_1 = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$   $f_1 = \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}$ 

are elements of  $\mathfrak{gl}_2$ .

For any smooth representation  $\rho : \operatorname{GL}_n \to \operatorname{GL}(V)$ , the differential of  $\rho$  is a Lie algebra representation  $\sigma : \mathfrak{gl}_n \to \mathfrak{gl}(V) = \operatorname{End}(V)$ . If  $v \in V$  is in the  $(p_1, \ldots, p_n)$  weight space, that is, if

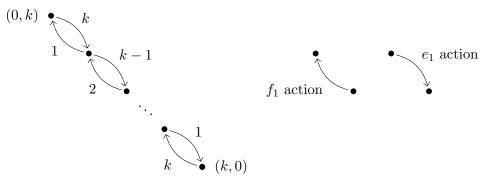
$$\rho \begin{pmatrix} t_1 & & \\ & \ddots & \\ & & t_n \end{pmatrix} \cdot v = (t_1^{p_1} \cdots t_n^{p_n})v,$$

then  $\sigma(E_{ij})v$  is in the  $(p_1, p_2, \ldots, p_i + 1, \ldots, p_j - 1, \ldots, p_n)$  weight space.

Consider the GL<sub>2</sub>-representation  $V_{(k)}(2)$  and the corresponding  $\mathfrak{gl}_2$ -representation. It has a weight basis with weights  $(k, 0), (k - 1, 1), \ldots, (0, k)$ . In that (ordered) basis, the  $\mathfrak{gl}_2$ -representation is given by

$$\sigma:\mathfrak{gl}_{2} \to \operatorname{End}(V_{(k)}(2)) \quad h_{1} \mapsto \begin{pmatrix} k & & & \\ & \ddots & & \\ & & 1 & \\ & & & 0 \end{pmatrix} \qquad h_{2} \mapsto \begin{pmatrix} 0 & & & \\ & 1 & & \\ & & \ddots & \\ & & & 0 \end{pmatrix} \qquad h_{2} \mapsto \begin{pmatrix} 0 & & & \\ & 1 & & \\ & & \ddots & \\ & & & k \end{pmatrix} \qquad e_{1} \mapsto \begin{pmatrix} 0 & & & & \\ & 0 & 2 & & \\ & & 0 & 2 & & \\ & & & 0 & & \\ & & & \ddots & k \\ & & & & 0 \end{pmatrix} \qquad f_{1} \mapsto \begin{pmatrix} 0 & & & & \\ & k & \ddots & & \\ & & \ddots & 0 & & \\ & & & & 1 & 0 \end{pmatrix}$$

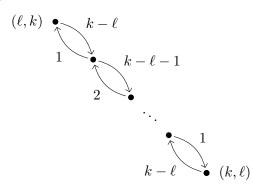
The actions of  $e_1$  and  $f_1$  on  $V_{(k)}(2)$  are captured in the following picture



#### AARON PRIBADI

where the dots (•) are weight spaces and the arrows for  $e_1$  and  $f_1$  are labeled with their factor.

For the representation  $V_{(k,\ell)}(2) \cong (\det)^{\otimes \ell} \otimes V_{(k-\ell)}$ , the picture for  $V_{(k-\ell)}$  is shifted to the weights  $(k,\ell), (k-1,\ell+1), \ldots, (\ell,k)$ .



**Remark:**  $\mathfrak{sl}_2$  controls the world!

# 2. Interaction of $\mathfrak{gl}_2$ strings

How do  $(e_j, f_j)$  act on  $(e_k, f_k)$  strings of a  $\mathfrak{gl}_n$ -representation?

**Proposition 1.** For any  $\mathfrak{gl}_n$ -representation:

- If  $|j k| \ge 2$ , then  $(e_j, f_j)$  preserves the length of the  $(e_k, f_k)$  string.
- If |j k| = 1, then  $e_j$  and  $f_j$  can only map between strings whose lengths differ by  $\pm 1$ .

*Proof.* Let  $(\sigma, W)$  be a  $\mathfrak{gl}_n$ -representation. Define

$$\phi_i:\mathfrak{gl}_2\to\mathfrak{gl}_n$$

by inserting the  $2 \times 2$  matrix in the rows *i* and i + 1 and the columns *i* and i + 1, with everything else zero.

Suppose that  $|j-k| \ge 2$ . Then  $[\phi_j(\mathfrak{gl}_2), \phi_k(\mathfrak{gl}_2)] = 0$ . Any two matrices  $u \in \phi_j(\mathfrak{gl}_2)$  and  $v \in \phi_k(\mathfrak{gl}_2)$  must commute, so  $\sigma(u) : W \to W$  is a map of  $\phi_k(\mathfrak{gl}_2)$ -representations. So the  $\phi_k(\mathfrak{gl}_2)$ -isotypic components of W are invariant under  $\sigma(u)$ , and  $\sigma(u)$  preserves the length of  $(e_k, f_k)$  strings.

The proof of the second statement uses the Serre relation, which will be introduced later.  $\Box$ 

**Example 2.** Consider  $V_{(2,1)}(3)$ . It has weight spaces

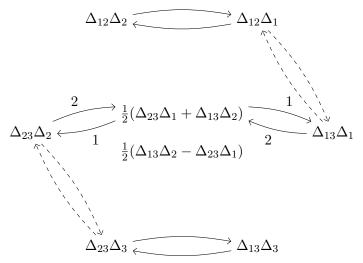
$$(1,2,0)$$
  $(2,1,0)$ 

$$(0,2,1)$$
  $(1,1,1)$   $(2,0,1)$ 

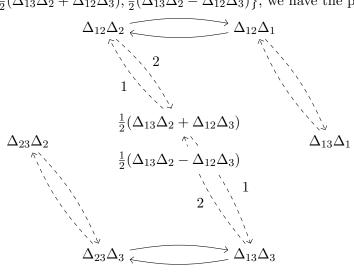
(0, 1, 2) (1, 0, 2)

and all of the weight spaces are one-dimensional except for (1, 1, 1), which is two-dimensional.

The (1,1,1) weight space contains the semistandard basis elements  $\{\Delta_{12}\Delta_3, \Delta_{13}\Delta_2\}$ . If we instead split the (1,1,1) weight space into the basis  $\{\frac{1}{2}(\Delta_{23}\Delta_1 + \Delta_{13}\Delta_2), \frac{1}{2}(\Delta_{23}\Delta_1 - \Delta_{13}\Delta_2\})$ , we have the picture



and with the basis  $\left\{\frac{1}{2}(\Delta_{13}\Delta_2 + \Delta_{12}\Delta_3), \frac{1}{2}(\Delta_{13}\Delta_2 - \Delta_{12}\Delta_3)\right\}$ , we have the picture



where the horizontal solid arrows are  $(e_1, f_1)$  strings and the diagonal dashed arrows are  $(e_2, f_2)$  strings.

The  $(e_1, f_1)$  and  $(e_2, f_2)$  strings come from the restriction of the GL<sub>3</sub>-representation  $V_{(2,1)}(3)$  to the subgroups

$$\begin{pmatrix} * & * & 0 \\ * & * & 0 \\ 0 & 0 & 1 \end{pmatrix} \quad \text{and} \quad \begin{pmatrix} 1 & 0 & 0 \\ 0 & * & * \\ 0 & * & * \end{pmatrix}$$

both isomorphic to  $GL_2$ .

For an example of the computations involved, consider the  $f_2$  string

$$\Delta_{12}\Delta_2 \xrightarrow{f_2} \frac{1}{2} (\Delta_{23}\Delta_1 + \Delta_{13}\Delta_2) \xrightarrow{f_2} \Delta_{13}\Delta_3.$$

Act on  $\Delta_{12}\Delta_2$  with the  $GL_3$  action by  $(1 + \epsilon f_2)$ 

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & \epsilon & 1 \end{pmatrix} \cdot \left( \begin{vmatrix} z_{11} & z_{12} \\ z_{11} & z_{12} \end{vmatrix} z_{12} \right) = \begin{vmatrix} z_{11} & z_{12} + \epsilon z_{13} \\ z_{21} & z_{22} + \epsilon z_{23} \end{vmatrix} (z_{12} + \epsilon z_{13})$$
$$= \Delta_{12}\Delta_2 + \epsilon (\Delta_{12}\Delta_3 + \Delta_{13}\Delta_2)$$

so that we get  $\Delta_{12}\Delta_3 + \Delta_{13}\Delta_2$ . The action of the Lie algebra also is a derivation (i.e. satisfies the Liebniz rule), so we can compute the same thing with

$$f_2(\Delta_{12}\Delta_2) = \Delta_{12}f_2(\Delta_2) + f_2(\Delta_{12})\Delta_2 = \Delta_{12}\Delta_3 + \Delta_{13}\Delta_2.$$

Note that it has the correct factor of 2 times the basis element.

The second segment of the string is

$$f_2(\frac{1}{2}(\Delta_{12}\Delta_3 + \Delta_{13}\Delta_2)) = \frac{1}{2}(f_2(\Delta_{12})\Delta_3 + \Delta_{12}f_2(\Delta_3) + f_2(\Delta_{13})\Delta_2 + \Delta_{13}f_2(\Delta_2))$$
  
=  $\frac{1}{2}(\Delta_{13}\Delta_3 + \Delta_{13}\Delta_3)$   
=  $\Delta_{13}\Delta_3$ 

as claimed.

From the pictures, one can also verify that the actions of  $(e_1, f_1)$  change the string length of  $(e_2, f_2)$  by  $\pm 1$ , and vice versa.

#### 3. Serre relation

**Proposition 3** (Serre relation). For any  $\mathfrak{gl}_n$ -representation  $\sigma : \mathfrak{gl}_n \to \operatorname{End}(W)$ , we have  $\sigma(f_k)^2 \sigma(f_{k+1}) - 2\sigma(f_k)\sigma(f_{k+1})\sigma(f_k) + \sigma(f_{k+1})\sigma(f_k)^2 = 0$ 

et cetera. Similar relations hold with  $f_k$  and  $f_{k+1}$  swapped, and with e's turned into f's.

For notation simplicity, we drop the  $\sigma$ 's.

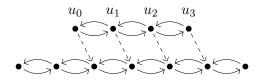
*Proof.* For the first equations,

$$\begin{aligned} f_k^2 f_{k+1} - 2f_k f_{k+1} f_k + f_{k+1} f_k^2 &= [f_k, f_k f_{k+1} - f_{k+1} f_k] \\ &= [f_k, [f_k, f_{k+1}]] \\ &= [E_{(k+1)k}, [E_{(k+1)k}, E_{(k+2)(k+1)}]] \\ &= [E_{(k+1)k}, -E_{(k+2)k}] = 0. \end{aligned}$$

The other equations are similar, and the map of Lie algebras preserves the bracket.

The Serre relation can be used to show that two  $(e_k, f_k)$  strings whose lengths differ by something other than  $\pm 1$  cannot be mapped to each other by  $f_{k+1}$  (or by  $e_{k+1}, f_{k-1}$ , or  $e_{k-1}$ ).

*Proof.* Suppose that the target string is at least two longer than the source string. Then at least one end has two extra nodes. Consider this illustration, where left arrows ( $\leftarrow$ ) are  $f_k$ , right arrows ( $\rightarrow$ ) are  $e_k$ , and dashed diagonal arrows ( $\searrow$ ) are  $f_{k+1}$ .

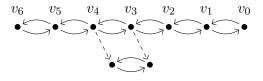


In this picture, in order to get the Serre relation

$$f_k^2 f_{k+1} - 2f_k f_{k+1} f_k + f_{k+1} f_k^2 = 0$$

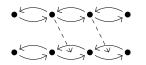
starting from the vector  $u_0$ , we must have  $f_{k+1} \cdot u_0$ . Similarly,  $f_{k+1} \cdot u_1 = 0$ , and by induction  $f_{k+1}u_i = 0$  for all  $u_i$  in the top string.

If the target string is at least two shorter than the source string, then we have this picture.



For the Serre relation to hold here,  $f_{k+1}(f_k^2(v_0)) = 0$ . Again by induction, all  $f_{k+1}$  maps are zero.

If the two strings are the same length, the  $f_{k+1}$  maps cannot go between the relevant weights.



The above arguments also work for  $f_{k+1}, e_{k+1}, f_{k-1}$ , and  $e_{k-1}$ . Thus, these maps can only map  $(e_k, f_k)$  strings to each other if the lengths of the strings differ by  $\pm 1$ .