


Dimers and Webs.

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Thomas Lam

University of Michigan

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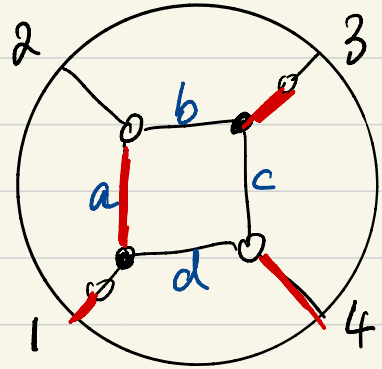
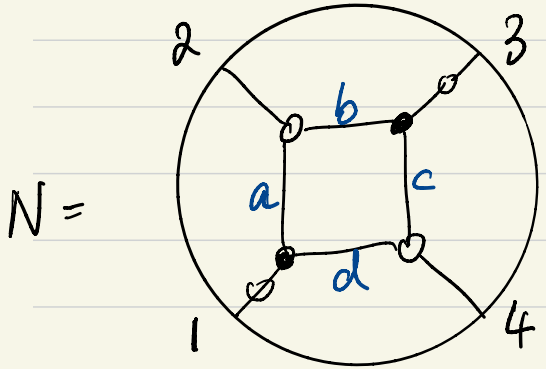
} Lecture 1

(j. C. Fraser and I. Le)

Lecture 2

Review: Dimer \rightarrow Grassmannian

(2)



$\partial(\pi) = \{1, 4\}$
 $wt(\pi) = a$

planar bipartite graph

boundary vertices are black and degree one

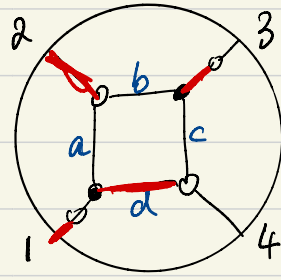
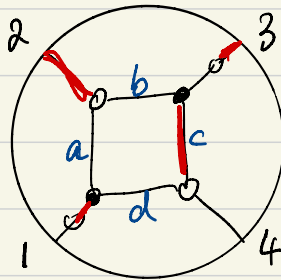
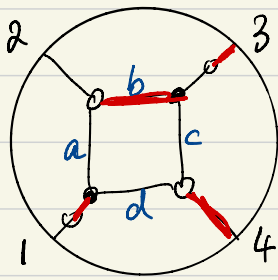
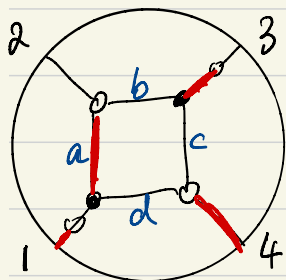
$\pi =$ almost perfect matching (uses all interior vertices) once.

$$\Delta_I(N) = \sum_{\pi: \partial(\pi) = I} wt(\pi)$$

$$|I| = k := \# \text{interior white} - \# \text{interior black}$$

Example of dimer generating function:

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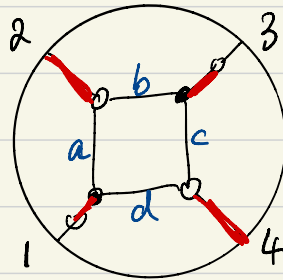
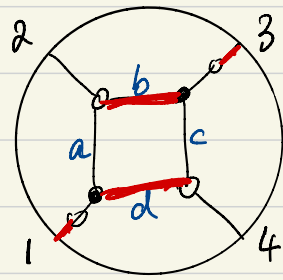
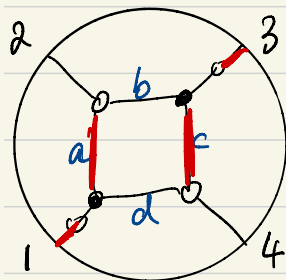
wt
∂

a
14

b
34

c
23

d
12



ac
13

bd
13

1
24

$$\Delta_{13}(N) = ac + bd$$

Theorem 1

[Kuo
Postnikov, Talaska, Postnikov - Speyer Williams
L]

$(\Delta_I(N))_{I \in \binom{[n]}{k}}$ defines a point in $Gr(k, n)$

ie. satisfies the Plücker relation.

$$\sum_{s=1}^{k+1} (-1)^s \Delta_{i_1, i_2, \dots, i_{k-1}, j_s} \Delta_{j_1, \dots, \hat{j}_s, \dots, j_{k+1}} = 0$$

e.g. $k=2$ $i_1=1$ $j_1, j_2, j_3 = 2, 3, 4$.

$$-\Delta_{12} \Delta_{34} + \Delta_{13} \Delta_{24} - \Delta_{14} \Delta_{23} = 0.$$

$$-d \cdot b + (act + bd) \cdot 1 - a \cdot c = 0$$

Proof using "webs"

Definition: A subgraph $\Sigma \subset N$ is 2-weblike if it is a union of connected components:

- (a) paths between boundary vertices
- (b) interior cycles
- (c) single edges (called "doubled edge")

such that every interior vertex is used.

$$F_{\mathcal{C}, T}(N) := \sum_{\Sigma} wt(\Sigma)$$

Temperley-Lieb Invariant.

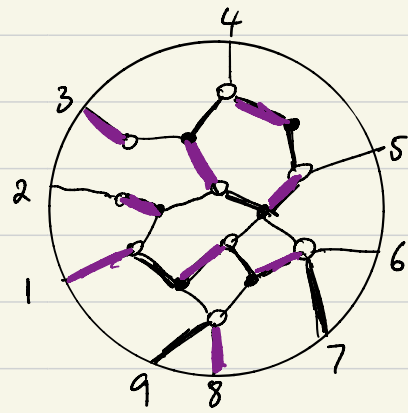
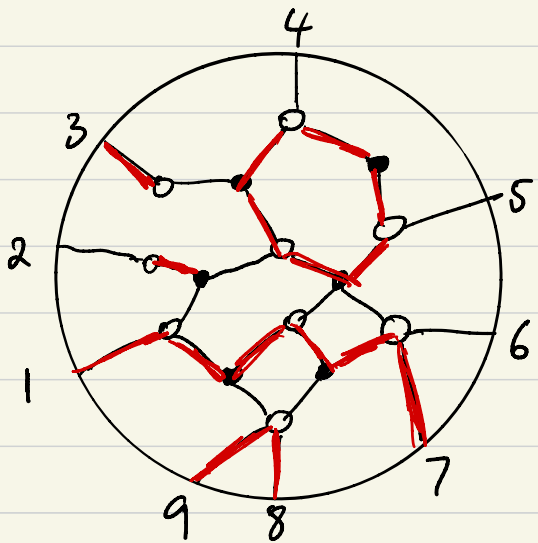
\mathcal{C} = partial non crossing matching

T = boundary vertices used in doubled edges

$$wt(\Sigma) = \prod_{\substack{e \in \text{path} \\ \text{or} \\ \text{cycle}}} wt(e) \cdot \prod_{\substack{e \in \\ \text{doubled} \\ \text{edge}}} wt(e)^2 \cdot 2^{\# \text{cycles}}$$

[L., cf. Rhoades -
Skandera
cf. Kenyon -
Wilson]

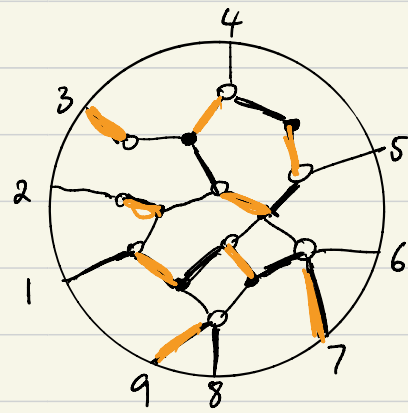
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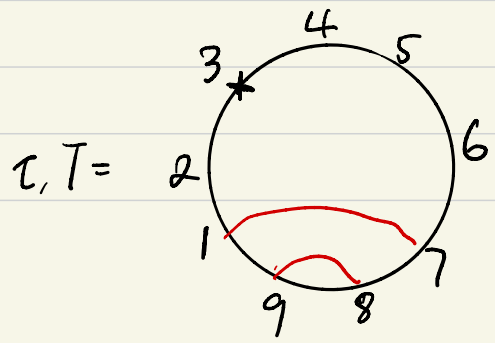
$$I = 138$$

Double
dimer

Σ



$$J = 379$$



(7)

Definition $I, J \in \binom{[n]}{k}$ is compatible with (τ, T) if.

- $T = I \cap J$
- each chord of τ matches $i \in I - J$ with $j \in J - I$

Theorem 2 [L. cf. Kenyon-Wilson, Rhoades-Skandera, ...]

$$\Delta_I(N) \Delta_J(N) = \sum_{\substack{(\tau, T) \\ \text{comp. with } (I, J)}} F_{\tau, T}(N)$$

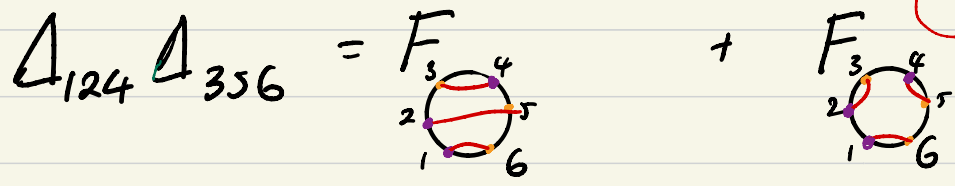
$$R(k, n) := \mathbb{C}[Gr(k, n)]$$

$$R(k, n) = \bigoplus_{d \geq 0} R(k, n)_d$$

$$R(k, n)_1 = \bigoplus \mathbb{C} \cdot \Delta_I$$

$$R(k, n)_2 = \bigoplus \mathbb{C} \cdot F_{\tau, T}$$

e.g.



- $F_{\tau, T}$ is a well-defined quadratic function on $Gr(\hat{k}, n)$
- forms a cyclically invariant basis
- coincides with "dual canonical basis", better than "standard monomial basis"

$$\Delta_{S_1} \Delta_{S_2} \quad \begin{array}{|c|c|} \hline S_1 & S_2 \\ \hline \end{array} \quad \begin{array}{l} \text{semi} \\ \text{standard} \end{array}$$

• Theorem 2 \Rightarrow Theorem 1

Expand $\sum_{s=1}^{k+1} (-1)^s \Delta_{i_1, i_2, \dots, i_{k-1}, j_s} \Delta_{j_1, \dots, j_s, \dots, j_{k+1}}$ in terms of $F_{\subseteq T}$

and do a sign-reversing involution.

• $\Delta_I \Delta_J \leq \Delta_{\min(I, J)} \Delta_{\max(I, J)} \leq \Delta_{\text{sort}_1(I, J)} \Delta_{\text{sort}_2(I, J)}$

e.g. $\Delta_{13567} \Delta_{23489} \leq \Delta_{13467} \Delta_{23589} \leq \Delta_{13468} \Delta_{23579}$
sort = 1233456789

[cf. Rhoades-Skandera
Farber-Postnikov
L.-Postnikov-Ryhyansky]

• $\mathcal{M}(N) := \{ I \in \binom{[n]}{k} \mid \Delta_I(N) \neq 0 \}$ matroid of N . ⑨

$I, J \in \mathcal{M} \Rightarrow \text{sort}_1, \text{sort}_2 \in \mathcal{M}$ [L.-Postnikov]
"sort-closed matroid"

• $\mathcal{M}^{(2)}(N) := \{ \tau, T \mid F_{\tau, T}(N) \neq 0 \}$ "quadratic matroid"

• any positive quadratic function, homogeneous with respect to torus action, is a positive

linear combination of $F_{\tau, T}$

• applications to dimers?

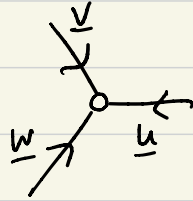
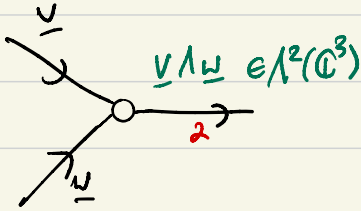
Webs

Based on joint work with C. Fraser and I. Le 10

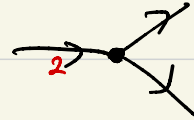
$$k \left\{ \left[\begin{array}{c} \downarrow v_1 \\ \downarrow v_2 \\ \dots \\ \downarrow v_n \end{array} \right] \right.$$

$$\mathbb{C}[Gr(k, n)] = R(k, n) = SL_k \text{ invariant functions on } \underbrace{\mathbb{C}^k \times \dots \times \mathbb{C}^k}_n = \mathbb{C}[\Delta_I] / \text{relations}$$

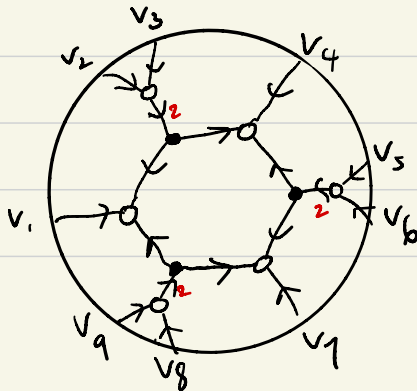
$k=3$



$$\Lambda^2 \mathbb{C}^3 \hookrightarrow \mathbb{C}^3 \otimes \mathbb{C}^3$$



$$\det \left(\frac{v}{w} \middle| \frac{u}{u} \right)$$



$$e \in R(3, 9)_3$$

$$v_i \in \mathbb{C}^3$$

[Reshetikhin-Turaev
Kuperberg]

Clusters and webs

(11)

[Fomin-Pylyavskyy]

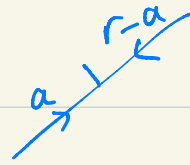
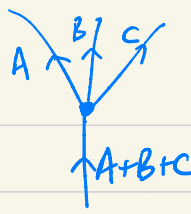
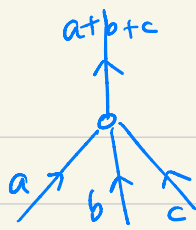
Cluster structure of $\mathbb{C}[Gr(3,n)]$ (+ more) can be described in terms of web invariants:

- initial cluster, quivers, coefficients

Conjecture }
[Fraser]
Gr(3,9)

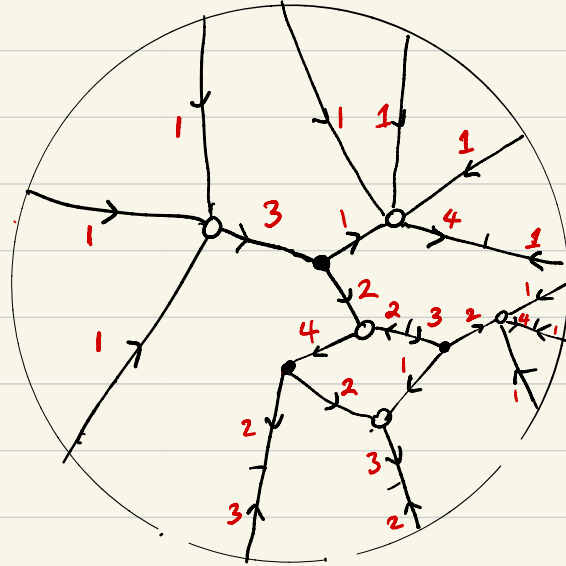
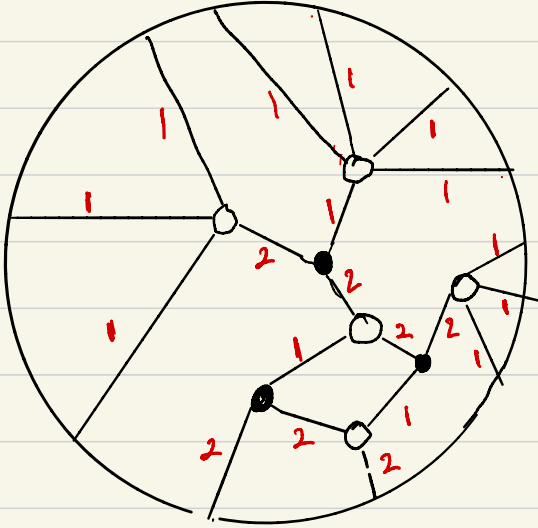
- cluster variables are indecomp. web invariants
- monomials \hookrightarrow tensor invariant
- compatibility

SL(r)-webs



(2)

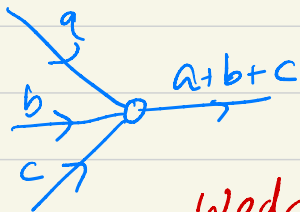
$r=5$



$$\vec{\lambda} = (1, 1, 1, \dots, 1, 2, 3)$$

untagged $SL(5)$ -web

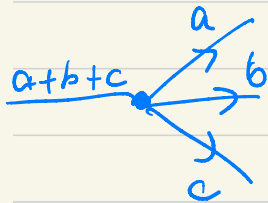
tagged $SL(5)$ -web



wedge

$$\Lambda^a(\mathbb{C}^r) \otimes \Lambda^b(\mathbb{C}^r) \otimes \Lambda^c(\mathbb{C}^r) \rightarrow \Lambda^{a+b+c}(\mathbb{C}^r)$$

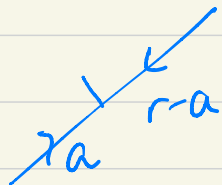
$$x_1 \otimes x_2 \otimes x_3 \mapsto x_1 \wedge x_2 \wedge x_3$$



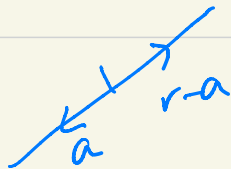
shuffle

$$\Lambda^{a+b+c}(\mathbb{C}^r) \rightarrow \Lambda^a(\mathbb{C}^r) \otimes \Lambda^b(\mathbb{C}^r) \otimes \Lambda^c(\mathbb{C}^r)$$

$$x_1 \wedge \dots \wedge x_{a+b+c} \rightarrow \sum \pm (x_{i_1} \wedge \dots \wedge x_{i_a}) \otimes (\quad) \otimes (\quad)$$



$$\Lambda^a(\mathbb{C}^r) \otimes \Lambda^{r-a}(\mathbb{C}^r) \rightarrow \Lambda^r(\mathbb{C}^r) \cong \mathbb{C}$$



$$\Lambda^r(\mathbb{C}^r) \rightarrow \Lambda^a(\mathbb{C}^r) \otimes \Lambda^{r-a}(\mathbb{C}^r)$$

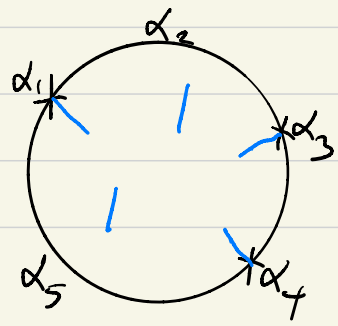
Lemma. Any tagging \hat{w} of a web W produces the same function up to sign.

$$W_\lambda(r) = \text{Hom}_{\text{SL}(r)} \left(\bigotimes_{i=1}^n \Lambda^{\lambda_i} \mathbb{C}^r, \mathbb{C} \right) \quad \lambda = (\lambda_1, \dots, \lambda_n)$$

$$W(r, n) := \text{Hom}_{\text{SL}(r)} \left(\bigotimes_{i=1}^n \mathbb{C}^r, \mathbb{C} \right)$$

are spanned by appropriate \hat{w} .

Ex. $r=1$

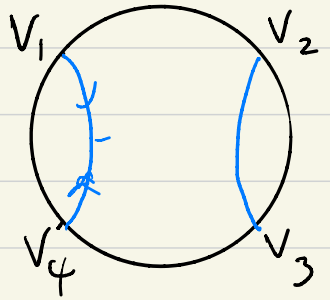


$\alpha_i \in \mathbb{C}$

$$\hat{w} = \pm \alpha_1 \alpha_3 \alpha_4$$

$r=2$

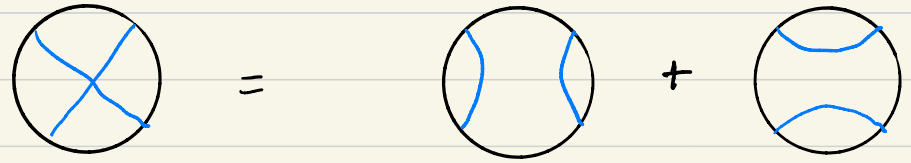
$v_i \in \mathbb{C}^2$



$$= \pm \det(v_1|v_4) \det(v_2|v_3)$$

basis of $W(2, 2n) =$ noncrossing matchings on $[2n]$

[Rumer-Teller
Wege

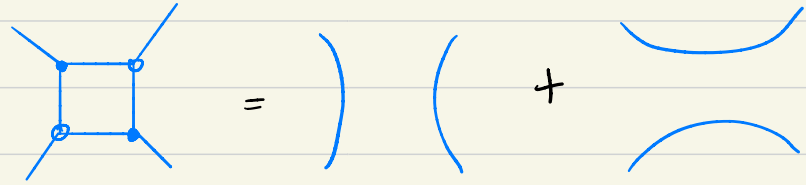
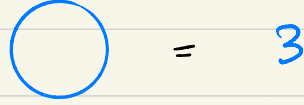
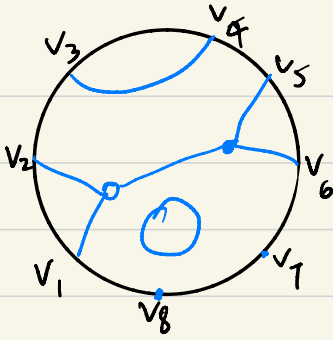


$$\text{circle} = \pm 2 = \text{trace}(\text{Id}: \mathbb{C}^2 \rightarrow \mathbb{C}^2)$$



$r=3$

e.g.



[Kuperberg] basis = irreducible/elliptic webs

ie. no contractible loops, no square faces

Definition An r -weblike subgraph $W \subset N$ is a subgraph using all vertices of N , with each edge labeled by a multiplicity $m(e)$, such that

$$\sum_{e \text{ incident to } v} m(e) = r \quad \text{for interior } v$$

$$\text{wt}(W) = \prod_e \text{wt}(e)^{m(e)}$$

$\lambda(W) =$ boundary multiplicities.

Each such W gives an untagged $SL(r)$ -web and a tensor invariant $\underline{W} (= \pm \hat{W}$ for some tagging)

$$\text{Web}_r(N; \lambda) := \sum_{\substack{\text{WCG} \\ \lambda(w) = \lambda}} \text{wt}(w) \underline{w} \in \mathcal{W}_\lambda(r)$$

• $r=1$ $\lambda \leftrightarrow I$ $\text{Web}_1(N; \lambda) = \Delta_I(N)$

Theorem 1
[Fraser, L., Le]

If $\Delta_I(N) = \Delta_I(N')$ for all I , then

$$\text{Web}_r(N; \lambda) = \text{Web}_r(N'; \lambda) \in \mathcal{W}_\lambda(r)$$

for all r, λ .

$\mathbb{C}[Gr(k,n)]_\lambda$ = be the λ -weight component.

e.g. $\Delta_{124} \Delta_{246}$ has weight $(1, 2, 0, 2, 0, 1)$

Theorem 2 Imm: $W_\lambda(U)^* \rightarrow \mathbb{C}[Gr(k,n)]_\lambda$

$$\varphi \mapsto (N \mapsto \varphi(\text{Web}_r(N; \lambda)))$$

is an isomorphism.

Theorem 3 $\Delta_{I_1}(N) \Delta_{I_2}(N) \dots \Delta_{I_r}(N) = \text{sign}(S) \text{Web}_r(N; \lambda) |_{e_S}$

where

$$S = (S(1), \dots, S(r)) \quad S(j) \subset [r]$$

I_i, S related by

$$I_i = \{j \in [n] : i \in S(j)\}$$

re. transpose as 01-matrices.

$$e_S = e_{S(1)} \otimes \dots \otimes e_{S(r)} \in \wedge^{d_1}(\mathbb{C}^r) \otimes \dots \otimes \wedge^{d_r}(\mathbb{C}^r)$$

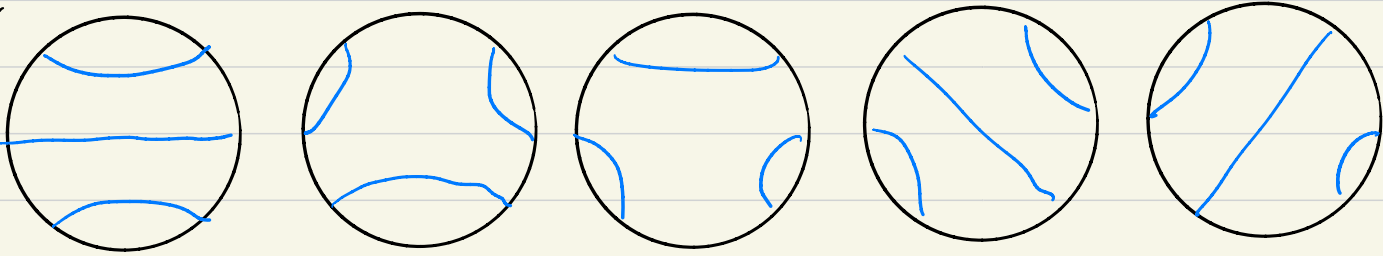
$$(e_{124} = e_1 \wedge e_2 \wedge e_4)$$

Remarks

• $r=2$

$W_\lambda(2)$ has noncrossing matching basis:

e.g.

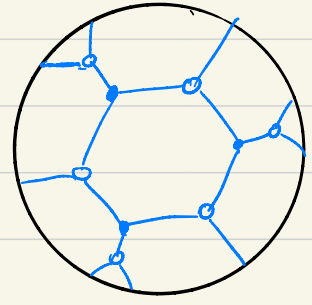
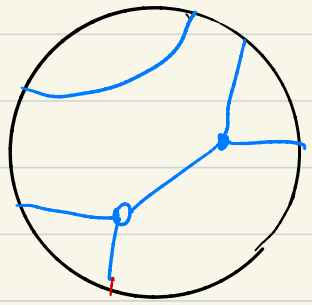


$$\in W_{(1,1,1,1,1)}(2) = W(2,6) \xleftrightarrow{\text{dual}} \mathbb{C}[Gr(3,6)]_{(1,1,1,1)}$$

$$\text{Imm}(\text{dual basis. element}) = F_{\tau, T}$$

• $r=3$

[Kuperberg] $W_\lambda(3)$ has basis of irreducible/elliptic webs.



no square faces!

$\text{Imm}(\text{dual basis element}) = \text{"web immanant"}$

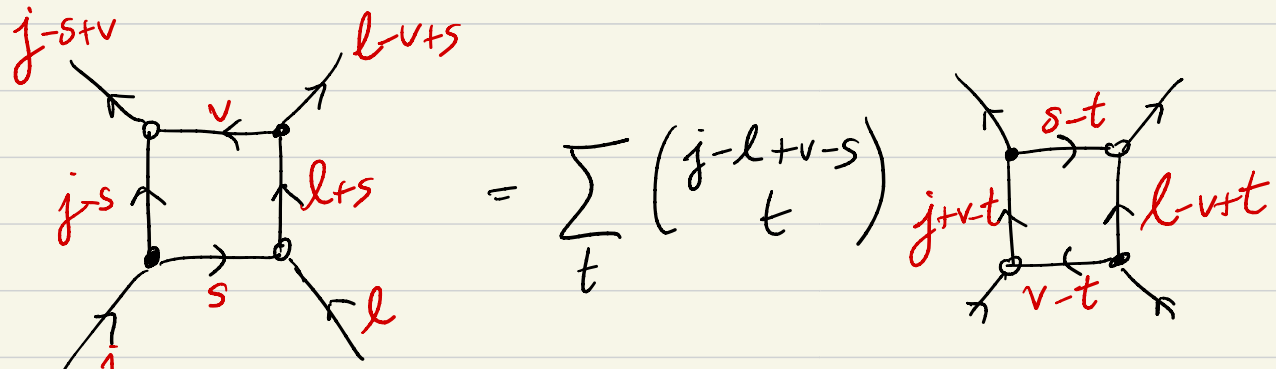
[L.
cf. Pylyavskyy]

- For $r > 3$, no generally agreed upon web basis is known.
There is a dual canonical basis.

generators for relations b/w webs are known
 [Kuperberg, ..., Cartier - Kamnitzer - Morrison]

Theorem Under $Web_r(N; \lambda) = Web_r(N'; \lambda)$

local moves for planar bipartite graphs \Rightarrow web relations



square move for tagged webs

Further directions

- Quantization?
- Positivity of \underline{W} as function on $Gr(r, n)$?
- Applications to dimers?
- Relation Imm to cluster structure?